# Galois connection for multiple-output operations 

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## Clones and coclones: the classical case

(1) Clones and coclones: the classical case

2 Interlude: reversible computing

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## Clones

Fix a base set $B$

## Definition

A clone is a set $\mathcal{C}$ of functions $f: B^{n} \rightarrow B, n \geq 0$, s.t.

- the projections $\pi_{n, i}: B^{n} \rightarrow B, \pi_{n, i}(\vec{x})=x_{i}$, are in $\mathcal{C}$
- $\mathcal{C}$ is closed under composition: if $g: B^{m} \rightarrow B$ and $f_{i}: B^{n} \rightarrow B$ are in $\mathcal{C}$, then

$$
h(\vec{x})=g\left(f_{0}(\vec{x}), \ldots, f_{m-1}(\vec{x})\right): B^{n} \rightarrow B
$$

is in $\mathcal{C}$

## Clones (cont'd)

- Clone generated by a set of functions $\mathcal{F}$
$=$ term functions of the algebra $(B, \mathcal{F})$
$=$ functions computable by circuits over $B$ using $\mathcal{F}$-gates
- Classical computing: clones on $B=\{0,1\}$ completely classified by [Post41]
- Clones can be studied by means of relations they preserve


## Preservation

$f: B^{n} \rightarrow B$ preserves $r \subseteq B^{k}:$

| $a_{0}^{0}$ | $\cdots$ | $a_{j}^{0}$ | $\cdots$ | $a_{n-1}^{0}$ |  | $b^{0}$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $\vdots$ |  | $\vdots$ |  | $\vdots$ | $f$ | $\vdots$ |  |
| $a_{0}^{i}$ | $\cdots$ | $a_{j}^{i}$ | $\cdots$ | $a_{n-1}^{i}$ |  | $b^{i}$ |  |
| $\vdots$ |  | $\vdots$ |  | $\vdots$ |  | $\vdots$ |  |
| $a_{0}^{k-1}$ | $\cdots$ | $a_{j}^{k-1}$ | $\cdots$ | $a_{n-1}^{k-1}$ |  | $b^{k-1}$ |  |
|  |  |  |  |  |  |  |  |
| $\pi$ | $\cdots$ | $\pi$ | $\cdots$ | $\pi$ |  | $\Longrightarrow$ | $\pi$ |
| $r$ |  | $r$ |  | $r$ |  | $r$ |  |

## Galois connection

$\mathcal{F}$ set of functions, $\mathcal{R}$ set of relations
Invariants and polymorphisms:

$$
\begin{aligned}
& \operatorname{Inv}(\mathcal{F})=\{r: \forall f \in \mathcal{F} f \text { preserves } r\} \\
& \operatorname{Pol}(\mathcal{R})=\{f: \forall r \in \mathcal{R} f \text { preserves } r\}
\end{aligned}
$$

$\Longrightarrow$ Galois connection: $\mathcal{R} \subseteq \operatorname{Inv}(\mathcal{F}) \Longleftrightarrow \mathcal{F} \subseteq \operatorname{Pol}(\mathcal{R})$

- $\operatorname{Pol}(\operatorname{Inv}(\mathcal{F})), \operatorname{Inv}(\operatorname{Pol}(\mathcal{R}))$ closure operators closed sets = range of Pol, Inv (resp.)
- Inv, Pol are mutually inverse dual isomorphisms of the complete lattices of closed sets


## Basic correspondence

## Theorem [Gei68,BKKR69]

If $B$ is finite:

- Galois-closed sets of functions = clones
- Galois-closed sets of relations = coclones


## Definition

Coclone $=$ set of relations closed under definitions by primitive positive FO formulas:

$$
R\left(x^{0}, \ldots, x^{k-1}\right) \Leftrightarrow \exists x^{k}, \ldots, x^{\prime} \bigwedge_{i<m} \varphi_{i}\left(x^{0}, \ldots, x^{\prime}\right)
$$

where each $\varphi_{i}$ is atomic

## Coclones (cont'd)

Equivalently: a set of relations is a coclone if it contains the identity $x_{0}=x_{1}$, and is closed under

- variable permutation and identification
- finite Cartesian products and intersections
- projection on a subset of variables

Closely related to constraint satisfaction problems

## Variants

A host of generalizations of this Galois connection appear in the literature (e.g., [Isk71,Ros71,Pös80,Ros83,Cou05,Ker12]):

- infinite base set
- partial functions, multifunctions
- functions $A^{n} \rightarrow B$
- categorial setting


# Interlude: reversible computing 

(1) Clones and coclones: the classical case
(2) Interlude: reversible computing

3 Clones and coclones revamped

## Computation in the physical world

Conventional models:
computation can destroy the input on a whim

$$
(x, y) \mapsto x+y
$$

Reality check:

## Landauer's principle

Erasure of $n$ bits of information incurs an $n k T \log 2$ increase of entropy elsewhere in the system
$\Longrightarrow$ dissipates energy as heat

Time-evolution operators in quantum mechanics are reversible

## Reversible computing

Reversible computation models:
only allow operations that can be inverted

$$
(x, y) \mapsto(x, x+y)
$$

Typical formalisms: circuits using reversible gates

- Classical computing:
- motivated by energy efficiency
- $n$-bit reversible gate $=$ permutation $\{0,1\}^{n} \rightarrow\{0,1\}^{n}$
- Quantum computing:
- $n$ qubits of memory $=$ Hilbert space $\mathbb{C}^{2^{n}}$
- quantum gate $=$ unitary linear operator
$\Longrightarrow$ inherently reversible


## Clones of reversible transformations

Reversible operations computable from a fixed set of gates:

- variable permutations, dummy variables
- composition
- ancilla bits: preset constant inputs, required to return to the original state at the end
$\Longrightarrow$ notion of "reversible clones"

Recently: [AGS15] gave complete classification for $B=\{0,1\}$
( $\approx$ Post's lattice for reversible operations)

# Clones and coclones revamped 

1 Clones and coclones: the classical case
(2) Interlude: reversible computing
(3) Clones and coclones revamped

## Goal

Generalize the clone-coclone Galois connection to encompass reversible clones

Let's first have a look at some simple reversible clones on $\{0,1\}$

## Examples

- Conservative operations $f:\{0,1\}^{n} \rightarrow\{0,1\}^{n}$ preserve Hamming weight

$$
f(\vec{a})=\vec{b} \Longrightarrow \sum_{i<n} a_{i}=\sum_{i<n} b_{i}
$$

- Mod-k preserving operations:

Hamming weight modulo $k$

$$
f(\vec{a})=\vec{b} \Longrightarrow \sum_{i<n} a_{i} \equiv \sum_{i<n} b_{i} \quad(\bmod k)
$$

Permutations "can count": invariants can't be just relations

## Examples (cont'd)

- Affine operations $f:\{0,1\}^{n} \rightarrow\{0,1\}^{n}$
$f(\vec{x})=A \vec{x}+\vec{c}$, where $\vec{c} \in \mathbb{F}_{2}^{n}, A \in \mathbb{F}_{2}^{n \times n}$ non-singular
- $\Longleftrightarrow$ each component $f_{i}:\{0,1\}^{n} \rightarrow\{0,1\}$ affine
- classical invariant: $f_{i}$ affine $\Longleftrightarrow$ preserves the relation $a+b+c+d=0$ on $\mathbb{F}_{2}^{4}$
- let $w: \mathbb{F}_{2}^{4} \rightarrow \mathbb{F}_{2}, w\left(a^{0}, a^{1}, a^{2}, a^{3}\right)=a^{0}+a^{1}+a^{2}+a^{3}$
- identify $\mathbb{F}_{2}=\{0,1\}=(\{0,1\}, 0, \vee)$
- $f:\{0,1\}^{n} \rightarrow\{0,1\}^{m}$ affine $\Longleftrightarrow$

$$
\begin{aligned}
& f\left(a_{0}^{0}, \ldots, a_{n-1}^{0}\right)=\left(b_{0}^{0}, \ldots, b_{m-1}^{0}\right), \ldots, \\
& f\left(a_{0}^{3}, \ldots, a_{n-1}^{3}\right)=\left(b_{0}^{3}, \ldots, b_{m-1}^{3}\right)
\end{aligned}
$$

implies

$$
\bigvee_{i<n} w\left(a_{i}^{0}, a_{i}^{1}, a_{i}^{2}, a_{i}^{3}\right) \geq \bigvee_{i<m} w\left(b_{i}^{0}, b_{i}^{1}, b_{i}^{2}, b_{i}^{3}\right)
$$

## General case

We consider a preservation relation between

- partial multifunctions $f: B^{n} \Rightarrow B^{m}$
- formally: $f \subseteq B^{n} \times B^{m}, n, m \geq 0$
- $f(\vec{x}) \approx \vec{y}$ denotes $(\vec{x}, \vec{y}) \in f$
- $\operatorname{Pmf}=\bigcup_{n, m} \operatorname{Pmf}_{n, m}$
- "weight functions" $w: B^{k} \rightarrow M$
- ( $M, 1, \cdot, \leq$ ) partially ordered monoid, $k \geq 0$
- $\mathrm{Wgt}=\bigcup_{k} \mathrm{Wgt}_{k}$


## Preservation

$f: B^{n} \Rightarrow B^{m}$ preserves $w: B^{k} \rightarrow M$ :

$$
\begin{array}{ccccccccc}
a_{0}^{0} & \cdots & a_{j}^{0} & \cdots & a_{n-1}^{0} & & b_{0}^{0} & \cdots & b_{m-1}^{0} \\
\vdots & & \vdots & & \vdots & f & \vdots & & \vdots \\
a_{0}^{i} & \cdots & a_{j}^{i} & \cdots & a_{n-1}^{i} & & & b_{0}^{i} & \cdots \\
\hline \vdots & & \vdots & & \vdots & & b_{m-1}^{i} \\
a_{0}^{k-1} \cdots & a_{j}^{k-1} & \cdots & a_{n-1}^{k-1} & & b_{0}^{k-1} \cdots & b_{m-1}^{k-1} \\
& & & & & & & & \\
& & & & & & \\
w\left(a_{0}\right) \cdots & w\left(a_{j}\right) & \cdots & w\left(a_{n-1}\right) & \leq & w\left(b_{0}\right) \cdots & \cdots\left(b_{m-1}\right)
\end{array}
$$

## Invariants and polymorphisms

The preservation relation induces a Galois connection

## Definition

$$
\text { If } \begin{aligned}
\mathcal{F} & \subseteq \operatorname{Pmf}, \mathcal{W} \subseteq \mathrm{Wgt}: \\
\operatorname{Inv}(\mathcal{F}) & =\{w \in \mathrm{Wgt}: \forall f \in \mathcal{F} f \text { preserves } w\} \\
\operatorname{Pol}(\mathcal{W}) & =\{f \in \operatorname{Pmf}: \forall w \in \mathcal{W} f \text { preserves } w\}
\end{aligned}
$$

What are the closed classes?

## Clones

$\operatorname{Pol}(\mathcal{W})$ has the following properties:

## Definition

$\mathcal{C} \subseteq \operatorname{Pmf}$ is a pmf clone if

- (identity) $\mathrm{id}_{n}: B^{n} \rightarrow B^{n}$ is in $\mathcal{C}$
- (composition) $\quad f: B^{n} \Rightarrow B^{m}, g: B^{m} \Rightarrow B^{r}$ in $\mathcal{C}$ $\Longrightarrow g \circ f: B^{n} \Rightarrow B^{r}$ in $\mathcal{C}$
- (products) $f: B^{n} \Rightarrow B^{m}, g: B^{n^{\prime}} \Rightarrow B^{m^{\prime}}$ in $\mathcal{C}$ $\Longrightarrow f \times g: B^{n+n^{\prime}} \Rightarrow B^{m+m^{\prime}}$ in $\mathcal{C}$

$$
(f \times g)\left(x, x^{\prime}\right) \approx\left(y, y^{\prime}\right) \Longleftrightarrow f(x) \approx y, g\left(x^{\prime}\right) \approx y^{\prime}
$$

- (topology) $\mathcal{C} \cap \mathrm{Pmf}_{n, m}$ is topologically closed $\ldots$


## Topological/local closure

Two interesting topologies on $\{0,1\}$ :

- $\mathbf{2}_{H}$ discrete (Hausdorff)
- $\mathbf{2}_{S}$ Sierpiński: $\{0\}$ closed, but $\{1\}$ not


## Lemma

Let $C \subseteq \mathcal{P}(X) \approx \mathbf{2}^{X}$. TFAE:

- $C$ is closed in $2_{S}^{X}$
- $C$ is closed in $2_{H}^{X}$ and under subsets
- $C$ is closed under directed unions and subsets
- $Y \in C$ iff all finite $Y^{\prime} \subseteq Y$ are in $C$

Previous slide: apply to $\operatorname{Pmf}_{n, m}=\mathcal{P}\left(B^{n} \times B^{m}\right)$

## Coclones

$\operatorname{Inv}(\mathcal{F})$ has the following properties:

## Definition

$\mathcal{D} \subseteq \mathrm{Wgt}$ is a weight coclone if

- (variable manipulation) $w: B^{k} \rightarrow M$ in $\mathcal{D}, \varrho: k \rightarrow 1$ $\Longrightarrow w\left(x^{\varrho(0)}, \ldots, x^{\varrho(k-1)}\right): B^{\prime} \rightarrow M$ in $\mathcal{D}$
- (homomorphisms) $w: B^{k} \rightarrow M$ in $\mathcal{D}, \varphi: M \rightarrow N$ $\Longrightarrow \varphi \circ \mathrm{w}: B^{k} \rightarrow N$ in $\mathcal{D}$
- (direct products) $\quad w_{\alpha}: B^{k} \rightarrow M_{\alpha}$ in $\mathcal{D} \quad(\alpha \in I)$ $\Longrightarrow\left(w_{\alpha}(x)\right)_{\alpha \in I}: B^{k} \rightarrow \prod_{\alpha \in I} M_{\alpha}$ in $\mathcal{D}$
- (submonoids) $w: B^{k} \rightarrow M$ in $\mathcal{D}, w\left[B^{k}\right] \subseteq N \subseteq M$ $\Longrightarrow w: B^{k} \rightarrow N$ in $\mathcal{D}$


## Galois connection

## Main theorem

For any $B$ :

- Galois-closed sets of pmf $=$ pmf clones
- Galois-closed classes of weights $=$ weight coclones


## Smaller invariants

Invariants of a pmf clone $\mathcal{C}$ form a proper class
Better: $\mathcal{C}=\operatorname{Pol}(\mathcal{W})$ s.t. for each $w: B^{k} \rightarrow M$ in $\mathcal{W}$ :

- $M$ is generated by $w\left[B^{k}\right]$
- call such weights tight
- $M$ finitely generated if $B$ finite
- $M$ is subdirectly irreducible (as a pomonoid)

Interesting case: (unordered) commutative monoids

- f.g. subdirectly irreducible are finite [Mal58]
- known structure [Sch66,Gri77]


## Variants

We might want to restrict Pmf or Wgt, or impose additional closure conditions, e.g.

- dimensions of $f: B^{n} \Rightarrow B^{m}$ :
- $n, m \geq 1, m=1, n=m$
- "kind" of $f$ :
- (partial/total) functions, permutations
- constraints on monoids:
- commutative, unordered
- constants, ancillas


## Monoid restrictions

- Classes of weights $w: B^{k} \rightarrow M$ with $M$ commutative $\Longleftrightarrow$ clones containing variable permutations

$$
\left(x_{0}, \ldots, x_{n-1}\right) \mapsto\left(x_{\pi(0)}, \ldots, x_{\pi(n-1)}\right)
$$

generated by swap $(x, y) \mapsto(y, x)$

- Classes of weights $w: B^{k} \rightarrow(M, 1, \cdot,=)$ (i.e., unordered monoids)
$\Longleftrightarrow$ clones closed under inverse

$$
f: B^{n} \Rightarrow B^{m} \text { in } \mathcal{C} \Longrightarrow f^{-1}: B^{m} \Rightarrow B^{n} \text { in } \mathcal{C}
$$

## Dimension constraints

$f: B^{n} \Rightarrow B^{m}$ with simple restrictions on $n, m$ form clones $\leftarrow$ ie $\Longrightarrow$ correspond to inclusion of particular weights:

- $n, m \geq 1$ : constant weight $c_{1}: B^{0} \rightarrow(2,1, \wedge,=)$
- $n=m: c_{1}: B^{0} \rightarrow(\mathbb{N}, 0,+,=)$
$m=1$ : clone $\mathcal{C}$ determined by $f: B^{n} \Rightarrow B$ iff contains swap \& diagonal maps $\Delta_{n}: B \rightarrow B^{n}, \Delta_{n}(x)=(x, \ldots, x)$

On the dual side:

- tight $w: B^{k} \rightarrow M$ in $\operatorname{Inv}(\mathcal{C})$ are $\{\wedge, \top\}$-semilattices
- subdirectly irreducible: $M=(2,1, \wedge, \leq)$
$\Longrightarrow$ weight functions $=$ relations
$\Longrightarrow$ agrees with the classical description


## Uniqueness conditions

Partial functions form a clone $\Longrightarrow$
$\mathcal{C}$ consists of partial functions iff
$\operatorname{Inv}(\mathcal{C})$ includes a particular weight:

- Kronecker delta $\delta: B^{2} \rightarrow(2,1, \wedge, \leq)$

Symmetrically:
$\mathcal{C}$ consists of injective pmf iff $\operatorname{lnv}(\mathcal{C})$ includes

$$
\delta: B^{2} \rightarrow(\{0,1\}, 1, \wedge, \geq)
$$

## Totality conditions

In the classical case:

- totality of functions in $\mathcal{C} \Longleftrightarrow$
closure of $\operatorname{Inv}(\mathcal{C})$ under existential quantification
- doesn't work well over infinite (uncountable) $B$


## Definition

$w: B^{k+1} \rightarrow(M, 1, \cdot, \leq)$ weight, $(M, 1, \cdot, 0,+)$ semiring
Define $w^{+}: B^{k} \rightarrow(M, 1, \cdot, \leq)$ by

$$
w^{+}\left(x^{0}, \ldots, x^{k-1}\right)=\sum_{u \in B} w\left(x^{0}, \ldots, x^{k-1}, u\right)
$$

## Orders on semirings

## Definition

- posemiring $=(M, 1, \cdot, 0,+, \leq)$ s.t.
- $(M, 1, \cdot, \leq)$ and $(M, 0,+, \leq)$ pomonoids
- ( $M, 1, \cdot, 0,+$ ) semiring
- positive semiring $=$ posemiring s.t. $0 \leq 1$ negative semiring $=$ posemiring s.t. $1 \leq 0$
- idempotent semiring: $x+x=x$ semilattice $\Longrightarrow$ can be ordered in two ways:
- $\vee$-semiring: + is $\vee$
$=$ idempotent positive semiring
- $\wedge$-semiring: + is $\wedge$
$=$ idempotent negative semiring


## Completeness of posemirings

## Definition

- complete idempotent semiring
( $\vee$-semiring, $\wedge$-semiring):
- complete lattice
- continuous idempotent semiring ( $\vee$-semiring, $\wedge$-semiring):
- complete
- infinite distributive laws

$$
\left(\sum_{i \in I} x_{i}\right) y=\sum_{i \in I} x_{i} y \quad y \sum_{i \in I} x_{i}=\sum_{i \in I} y x_{i}
$$

Continuous $\vee$-semirings $=$ unital quantales

## Total clones

$\mathcal{C}=\operatorname{Pol}(\mathcal{D}), \mathcal{D}=\operatorname{Inv}(\mathcal{C})$
For $B$ countable, the following are equivalent:

- $\mathcal{C}$ is generated by total multifunctions
- $w: B^{k+1} \rightarrow M$ is in $\mathcal{D}, M$ is a continuous $\vee$-semiring $\Longrightarrow w^{+}: B^{k} \rightarrow M$ is in $\mathcal{D}$

Symmetrically: clones of surjective pmf characterized using continuous $\wedge$-semirings

For $B$ finite, TFAE:

- $\mathcal{C}$ is generated by mf extending a bijective function
- $w: B^{k+1} \rightarrow M$ is in $\mathcal{D}, M$ is a posemiring

$$
\Longrightarrow w^{+}: B^{k} \rightarrow M \text { is in } \mathcal{D}
$$

## Ancillas

$$
\mathcal{C}=\operatorname{Pol}(\mathcal{D}), \mathcal{D}=\operatorname{lnv}(\mathcal{C})
$$

The following are equivalent:

- $\mathcal{C}$ supports ancillas
$c \in B, f: B^{n+1} \Rightarrow B^{m+1}$ in $\mathcal{C} \Longrightarrow f_{c}: B^{n} \Rightarrow B^{m}$ in $\mathcal{C}$

$$
f_{c}(\vec{x}) \approx \vec{y} \Longleftrightarrow f(\vec{x}, c) \approx(\vec{y}, c)
$$

- $\mathcal{D}$ is generated by $w: B^{k} \rightarrow M$ s.t. the diagonal weights $z=w(u, \ldots, u)$ for $u \in B$ are right-order-cancellative

$$
x z \leq y z \Longrightarrow x \leq y
$$

Warning: Interferes badly with totality

## Summary

- The standard clone-coclone duality extends to a Galois connection between partial multifunctions $B^{n} \Rightarrow B^{m}$ and pomonoid-valued functions $B^{k} \rightarrow M$
- Gracefully restricts to natural subclasses, such as total functions $B^{n} \rightarrow B^{m}$


## Thank you for attention!

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