

# Elementary analytic functions in $VTC^0$

Emil Jeřábek

Institute of Mathematics  
Czech Academy of Sciences  
jerabek@math.cas.cz  
<http://math.cas.cz/~jerabek/>

Czech Gathering of Logicians  
June 2022, Prague

# $TC^0$ and $VTC^0$

- 1**  $TC^0$  and  $VTC^0$
- 2 Iterated multiplication and division
- 3 Induction in  $VTC^0$
- 4 Elementary analytic functions

# Theories vs. complexity classes

Correspondence of theories of bounded arithmetic  $T$  and computational complexity classes  $C$ :

- ▶ Provably total computable functions of  $T$  are  $C$ -functions
- ▶  $T$  can reason using  $C$ -predicates (comprehension, induction, minimization, ...)

Feasible reasoning:

- ▶ Given a concept  $X \in C$ , what can we prove about  $X$  while reasoning only with concepts from  $C$ ?
- ▶ Formalization: what does  $T$  prove about  $X$ ?

This talk:

$X$  = elementary integer arithmetic operations  $+$ ,  $\cdot$ ,  $\leq$

# The class $TC^0$

$$AC^0 \subseteq ACC^0 \subseteq TC^0 \subseteq NC^1 \subseteq L \subseteq NL \subseteq AC^1 \subseteq \dots \subseteq P$$

- $TC^0$  = dlogtime-uniform  $O(1)$ -depth  $n^{O(1)}$ -size  
unbounded fan-in circuits with threshold gates  
= **FOM**-definable on finite structures  
representing strings  
(first-order logic with majority quantifiers)  
=  $O(\log n)$  time,  $O(1)$  thresholds  
on a threshold Turing machine  
= Constable's  $\mathcal{K}$ : closure of  $+, -, \cdot, /$  under substitution  
and polynomially bounded  $\Sigma, \Pi$

# $\mathbf{TC}^0$ and arithmetic operations

For integers given in binary:

- ▶  $+$  and  $\leq$  are in  $\mathbf{AC}^0 \subseteq \mathbf{TC}^0$
- ▶  $\times$  is in  $\mathbf{TC}^0$  ( $\mathbf{TC}^0$ -complete under  $\mathbf{AC}^0$  reductions)

$\mathbf{TC}^0$  can also do:

- ▶ iterated addition  $\sum_{i < n} X_i$
- ▶ integer division and iterated multiplication  
[BCH'86,CDL'01,HAB'02]
- ▶ the corresponding operations on  $\mathbb{Q}$ ,  $\mathbb{Q}(\alpha)$ , ...
- ▶ approximate functions given by nice power series:
  - ▶  $\sin X$ ,  $\log X$ ,  $\sqrt[k]{X}$ , ...
- ▶ sorting, ...

$\implies \mathbf{TC}^0$  is the right class for basic arithmetic operations

# The theory $VTC^0$

- ▶ Zambella-style **two-sorted** bounded arithmetic
  - ▶ unary (auxiliary) integers with  $0, 1, +, \cdot, \leq$
  - ▶ finite sets = binary integers = binary strings
- ▶ Noteworthy axioms:
  - ▶  $\Sigma_0^B$ -comprehension ( $\Sigma_0^B =$  bounded, w/o SO q'ifiers)
  - ▶ every set has a counting function
- ▶ Correspondence to  $\mathbf{TC}^0$ :
  - ▶ provably total computable (i.e.,  $\exists \Sigma_0^B$ -definable) functions are exactly the  $\mathbf{TC}^0$ -functions
  - ▶ has induction, minimization, ... for  $\mathbf{TC}^0$ -predicates
- ▶ Basic binary integer arithmetic in  $VTC^0$ :
  - ▶ can define  $+, \cdot, \leq$  on binary integers
  - ▶ proves integers form a discretely ordered ring (*DOR*)

# $TC^0$ feasible reasoning

What else can  $VTC^0$  do with basic arithmetic operations?

- ▶ [J'22] **Iterated multiplication and division:**  
formalize a variant of the [HAB'02] algorithm
- ▶ [J'15] **Open induction** in  $\langle +, \cdot, < \rangle$  (*IOpen*),  
(translation of)  $\Sigma_0^b$ -minimization in Buss's language
- ▶ [J'??] **Elementary analytic functions:**  
exp, log, sin, arcsin, sinh, arsinh, ...

# Iterated multiplication and division

- 1  $TC^0$  and  $VTC^0$
- 2 Iterated multiplication and division**
- 3 Induction in  $VTC^0$
- 4 Elementary analytic functions



# History

[BCH'86]

- ▶  $\prod_{i < n} X_i$ ,  $\lfloor Y/X \rfloor$ ,  $X^n$  are  $\mathbf{TC}^0$ -reducible to each other
- ▶ they are in  $\mathbf{P}$ -uniform  $\mathbf{TC}^0$
- ▶ compute the product in Chinese remainder representation:

$$\text{CRR}_{\vec{m}}(X) = \langle X \bmod m_i : i < k \rangle$$

where  $\vec{m} = \langle m_i : i < k \rangle$  small primes

- ▶ (NB: predates definition of  $\mathbf{TC}^0$ )

Improved CRR reconstruction procedures  $\implies$

- ▶ [CDL'01]: logspace-uniform  $\mathbf{TC}^0$  (hence  $\mathbf{L}$ )
- ▶ [HAB'02]: dlogtime-uniform  $\mathbf{TC}^0$

# Formalization in $VTC^0$

Raised as a problem by Atserias [Ats'03,NC'06]

## Obstacles:

- ▶ complex structure with interdependent parts
- ▶ analysis elementary, but **chicken-and-egg** problems:  
uses iterated products and divisions all over the place

## Results [J'22]:

- ▶  $VTC^0$  proves *IMUL* and *DIV*
- ▶  $I\Delta_0 + WPHP(\Delta_0)$  has a well-behaved  $\Delta_0$  definition of  $a^r \bmod m$

## Induction in $VTC^0$

- 1  $TC^0$  and  $VTC^0$
- 2 Iterated multiplication and division
- 3 Induction in  $VTC^0$**
- 4 Elementary analytic functions

# Open induction

**Question:** Can  $VTC^0$  prove some amount of induction for binary numbers?

The weakest nontrivial fragment of induction:  $IOpen$

- ▶ induction for **quantifier-free** formulas in language  $\langle +, \cdot, < \rangle$
- ▶ [Shep'64]  $\mathfrak{M} \models IOpen \iff$   
 $\mathfrak{M}$  is an **integer part** of a **real-closed field**

$VTC^0$ -provable  $\forall\exists\Sigma_0^B$  statements witnessed by  $\mathbf{TC}^0$  functions  
 $\implies$  the following are equivalent:

- ▶  $VTC^0 \vdash IOpen$
- ▶ for every constant  $d$ ,  $VTC^0$  can formalize a  $\mathbf{TC}^0$  **root approximation algorithm** for degree- $d$  polynomials

# Results [J'15]

$VTC^0$  does prove *IOpen*:

- ▶ **Largange inversion formula**  $\implies$  approximation of roots of polynomials with “small” constant coefficient
- ▶ model-theoretic argument using Shepherdson's criterion
  - ▶  $\mathfrak{M} \rightsquigarrow \text{DOR } \mathbf{Z}^{\mathfrak{M}} \rightsquigarrow$  fraction field  $\mathbf{Q}^{\mathfrak{M}} \rightsquigarrow$  completion  $\mathbf{R}^{\mathfrak{M}}$
  - ▶  $\mathfrak{M} \models \text{DIV} \implies \mathbf{Z}^{\mathfrak{M}}$  integer part of  $\mathbf{Q}^{\mathfrak{M}}$  and  $\mathbf{R}^{\mathfrak{M}}$
  - ▶ LIF  $\implies \mathbf{R}^{\mathfrak{M}}$  henselian  $\implies \mathbf{R}^{\mathfrak{M}}$  real-closed

Extend the argument using ideas of [Man'91]:

- ▶  $VTC^0$  proves **induction** and **minimization** for translations of  $\Sigma_0^b$  formulas in Buss's language

# Elementary analytic functions

- 1  $TC^0$  and  $VTC^0$
- 2 Iterated multiplication and division
- 3 Induction in  $VTC^0$
- 4 **Elementary analytic functions**

# $\mathbf{TC}^0$ analytic functions

Recall:  $\mathbf{TC}^0$  can compute approximations of analytic functions whose power series have  $\mathbf{TC}^0$ -computable coefficients

Question: Can  $VTC^0$  prove their basic properties?

There's a plethora of such functions  $\implies$  let's start small:

Elementary analytic functions (real and complex)

- ▶ exp, log
- ▶ trigonometric: sin, cos, tan, cot, sec, csc
- ▶ inverse trig.: arcsin, arccos, arctan, arccot, arcsec, arccsc
- ▶ hyperbolic: sinh, cosh, tanh, coth, sech, csch
- ▶ inverse hyp.: arsinh, arcosh, artanh, arcoth, arsech, arcsch

All definable in terms of complex exp and log

# $VTC^0$ setup

Working with rational approximations only is **quite tiresome**

Recall:  $\mathfrak{M} \models VTC^0 \rightsquigarrow \text{DOR } \mathbf{Z}^{\mathfrak{M}} \rightsquigarrow \text{fraction field } \mathbf{Q}^{\mathfrak{M}}$   
 $\rightsquigarrow \text{completion } \mathbf{R}^{\mathfrak{M}} \rightsquigarrow \mathbf{C}^{\mathfrak{M}} = \mathbf{R}^{\mathfrak{M}}(i)$

Treat the functions as  $f: \mathbf{C}^{\mathfrak{M}} \rightarrow \mathbf{C}^{\mathfrak{M}}$  (or on a subset)

This simplifies development, but approximations **still needed**:

- ▶ translate results back to the language of  $VTC^0$
- ▶ use the functions in induction arguments, ...

Further notation: unary integers embed as  $\mathbf{L}^{\mathfrak{M}} \subseteq \mathbf{Z}^{\mathfrak{M}}$

$\mathbf{C}_L^{\mathfrak{M}} = \{z \in \mathbf{C}^{\mathfrak{M}} : \exists n \in \mathbf{L}^{\mathfrak{M}} |z| \leq n\}$ ,  $\mathbf{R}_L^{\mathfrak{M}} = \mathbf{R}^{\mathfrak{M}} \cap \mathbf{C}_L^{\mathfrak{M}}$ , ...



# Main results

We can define  $\pi \in \mathbf{R}^m$ ,

$$\exp: \mathbf{R}_L^m + i\mathbf{R}^m \rightarrow \mathbf{C}_{\neq 0}^m,$$

$$\log: \mathbf{C}_{\neq 0}^m \rightarrow \mathbf{R}_L^m + i(-\pi, \pi]$$

such that

- ▶  $\exp(z_0 + z_1) = \exp z_0 \exp z_1$
- ▶  $\exp$  is  $2\pi i$ -periodic
- ▶  $\exp \log z = z$
- ▶  $\log \exp z = z$  for  $z \in \mathbf{R}_L^m + i(-\pi, \pi]$
- ▶  $\exp \upharpoonright \mathbf{R}_L^m$  increasing bijection  $\mathbf{R}_L^m \rightarrow \mathbf{R}_{>0}^m$ , convex
- ▶ for small  $z$ :  $\exp z = 1 + z + O(z^2)$ ,  $\log(1 + z) = z + O(z^2)$

# Outline of the construction

- ▶ Define  $\exp: \mathbf{C}_{\mathbb{L}}^{\mathfrak{m}} \rightarrow \mathbf{C}^{\mathfrak{m}}$  using  $\sum_n \frac{z^n}{n!}$   
show  $\exp(z_0 + z_1) = \exp z_0 \exp z_1$
- ▶ Define  $\log$  on a nbh of 1 using  $-\sum_n \frac{(1-z)^n}{n}$   
show  $\log(z_0 z_1) = \log z_0 + \log z_1$  for  $z_j$  close enough to 1
- ▶ Extend  $\log$ 
  - ▶ to  $\mathbf{R}_{>0}^{\mathfrak{m}}$  using  $2^n: \mathbf{L}^{\mathfrak{m}} \rightarrow \mathbf{Z}^{\mathfrak{m}}$
  - ▶ to an angular sector by combining the two
  - ▶ to  $\mathbf{C}_{\neq 0}^{\mathfrak{m}}$  using  $8 \log \sqrt[8]{z}$
- ▶  $\log \exp(z_0 + z_1) = \log \exp z_0 + \log \exp z_1$  when  $|\operatorname{Im} z_j|$  small  
 $\implies \log \exp z = z$  when  $|\operatorname{Im} z|$  small  
 $\implies \exp \log z = z$  using injectivity of  $\log$
- ▶  $\exp$  is  $2\pi i$ -periodic for  $\pi := \operatorname{Im} \log(-1)$   
 $\implies$  extend  $\exp$  to  $\mathbf{R}_{\mathbb{L}}^{\mathfrak{m}} + i\mathbf{R}^{\mathfrak{m}}$

# Applications

Define

- ▶  $z^w = \exp(w \log z)$ ,  $\sqrt[n]{z} = z^{1/n}$
- ▶  $\prod_{j < n} z_j$  for a sequence of  $z_j \in \mathbf{Q}^{\mathfrak{M}}(i)$  coded in  $\mathfrak{M}$
- ▶ trigonometric, inverse trigonometric, hyperbolic, inverse hyperbolic functions

Model-theoretic consequence:

- ▶ Every countable model of  $VTC^0$  is an exponential integer part of a real-closed exponential field (even though  $\exp$  is not total on  $\mathbf{R}^{\mathfrak{M}}$  !)

# References

- ▶ A. Atserias: Improved bounds on the Weak Pigeonhole Principle and infinitely many primes from weaker axioms, Theoret. Comput. Sci. 295 (2003), 27–39
- ▶ P. Beame, S. Cook, H. Hoover: Log depth circuits for division and related problems, SIAM J. Comp. 15 (1986), 994–1003
- ▶ A. Chiu, G. Davida, B. Litow: Division in logspace-uniform  $\text{NC}^1$ , RAIRO – Theoret. Inf. Appl. 35 (2001), 259–275
- ▶ S. Cook, P. Nguyen: Logical foundations of proof complexity, Cambridge Univ. Press, 2010
- ▶ W. Hesse, E. Allender, D. M. Barrington: Uniform constant-depth threshold circuits for division and iterated multiplication, J. Comp. System Sci. 65 (2002), 695–716

# References (cont'd)

- ▶ E. Jeřábek: *Open induction in a bounded arithmetic for  $TC^0$* , Arch. Math. Logic 54 (2015), 359–394
- ▶ E. Jeřábek: *Iterated multiplication in  $VTC^0$* , Arch. Math. Logic (2022), <https://doi.org/10.1007/s00153-021-00810-6>
- ▶ E. Jeřábek: *Elementary analytic functions in  $VTC^0$* , in preparation
- ▶ E. Jeřábek: *Models of  $VTC^0$  as exponential integer parts, ?*
- ▶ S.-G. Mantzivis: *Circuits in bounded arithmetic part I*, Ann. Math. Artif. Intel. 6 (1991), 127–156
- ▶ P. Nguyen, S. Cook: *Theories for  $TC^0$  and other small complexity classes*, Log. Methods Comput. Sci. 2 (2006), art. 3
- ▶ J. Shepherdson: *A nonstandard model for a free variable fragment of number theory*, Bull. Acad. Polon. Sci. 12 (1964), 79–86